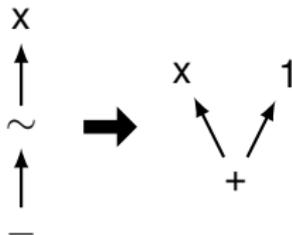
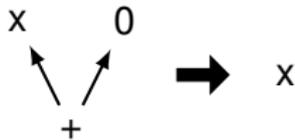


OPTGEN: A Generator for Local Optimizations

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Local optimizations:

- IR level
 - SSA form
 - Data dependency graph
- Do not require any global analysis
- Can be applied at any time during compilation

Goal

Generate all local optimizations (up to a given cost limit).

Input:

- Set of operations and their costs
- Cost limit
- Bit width

Output:

- Complete set of verified local optimizations

Assembly level

```
mov x, r0  
mov y, r1
```

```
xor r0, r1, r2  
...  
...  
or r0, r2, r3
```

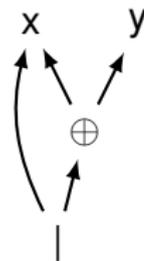


```
mov x, r0  
mov y, r1
```

```
...  
...  
or r0, r1, r3
```

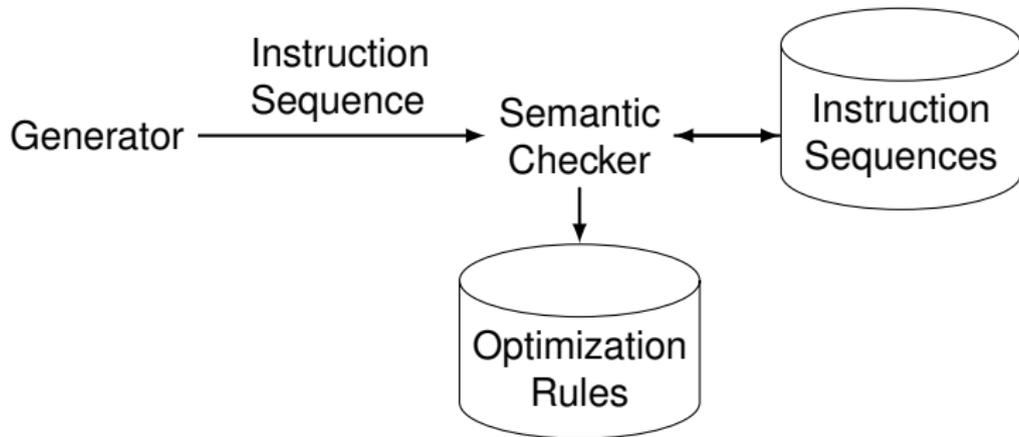
- Peephole of k instructions
- Architecture-specific
- Precise cost model

IR level



- Pattern of k values
- Independent of Architecture
- SSA form

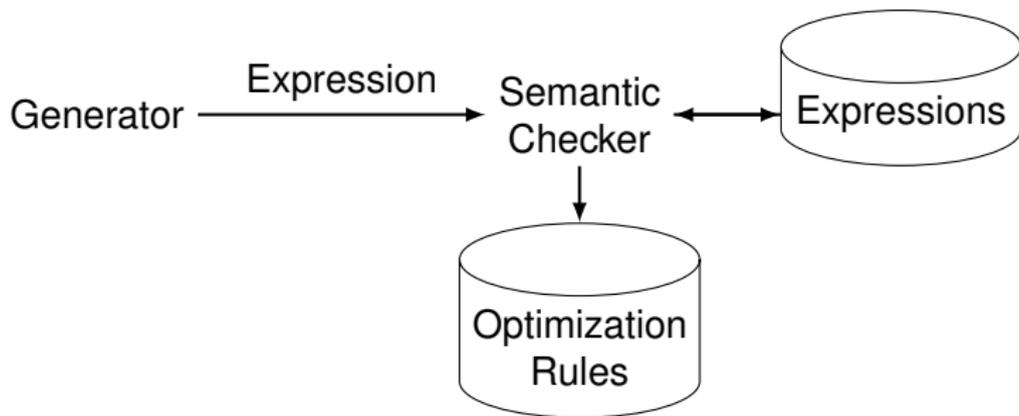
Common Design of Peephole Generators



Generator Generates all possible instructions sequences

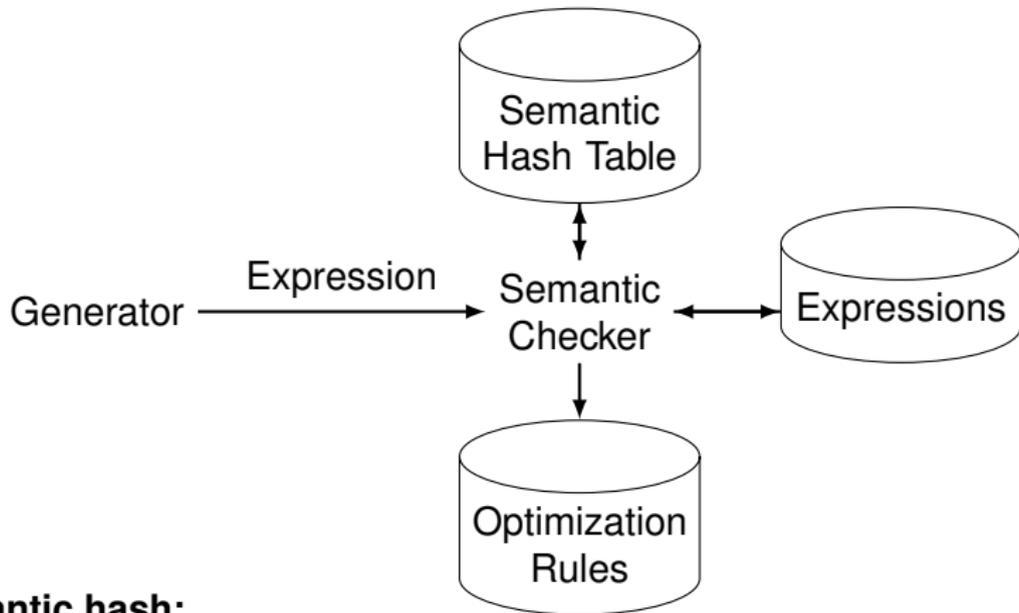
Semantic Checker Proofs the equivalence of two instruction sequences

Design of OPTGEN (so far)



Generator Generates all possible expressions

Semantic Checker Proofs the equivalence of two expressions



Semantic hash:

- Evaluate expression for precomputed test inputs
- $\text{semantic_hash}(x) = \text{semantic_hash}(x \mid 0)$

OPTGEN **parameters:**

- Operations:
 - Constants (cost: 0)
 - And (cost: 1)
 - Or (cost: 1)
 - Not (cost: 1)
- Cost limit: 2
- Bit width: 8

Example – Costs 0

Enumerate expressions with costs 0:

- x
- 0
- 1
- ...
- 255

Combine expressions with existing operations:

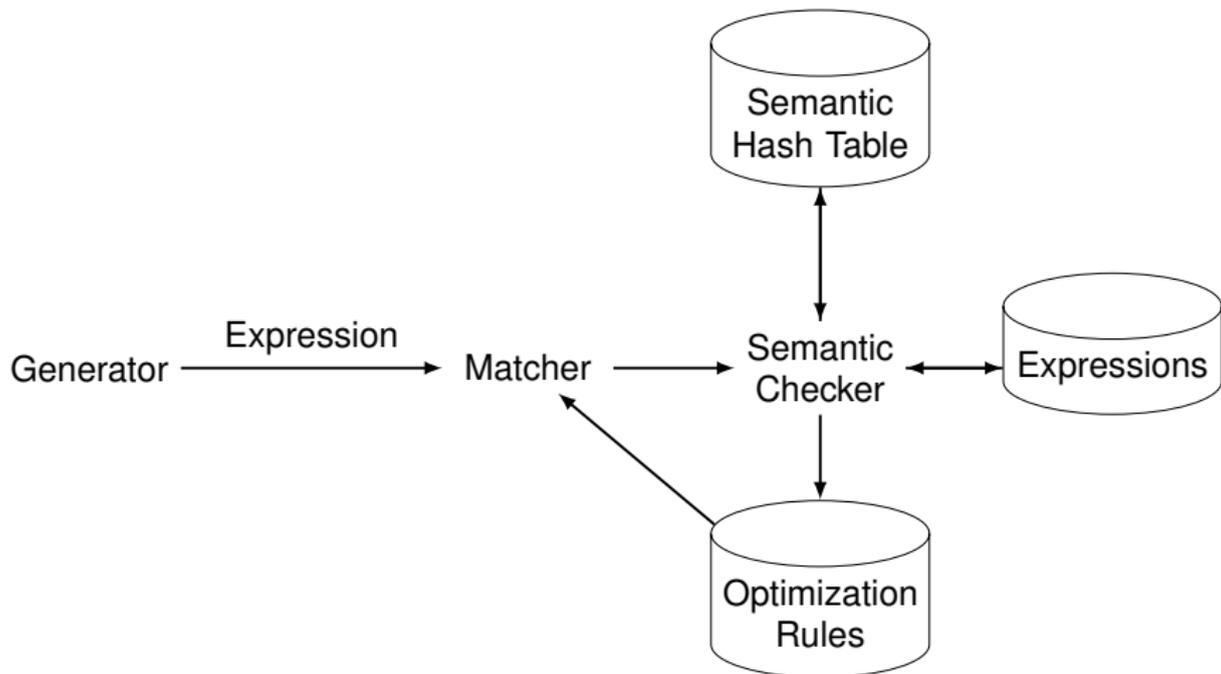
- y
- $x \& x$
 - Same semantic hash class as x
 - SMT check: $x \& x = x$
 - Optimization: $x \& x \rightarrow x$
- $x \& 0$
 - Same semantic hash class as 0
 - SMT check: $x \& 0 = 0$
 - Optimization: $x \& 0 \rightarrow 0$

Example – Costs 2

Combine expressions with existing operations:

- $(x \& y) \& 0$
 - Rule $x \& 0 \rightarrow 0$ applicable
 - No further action

Design of OPTGEN (so far)



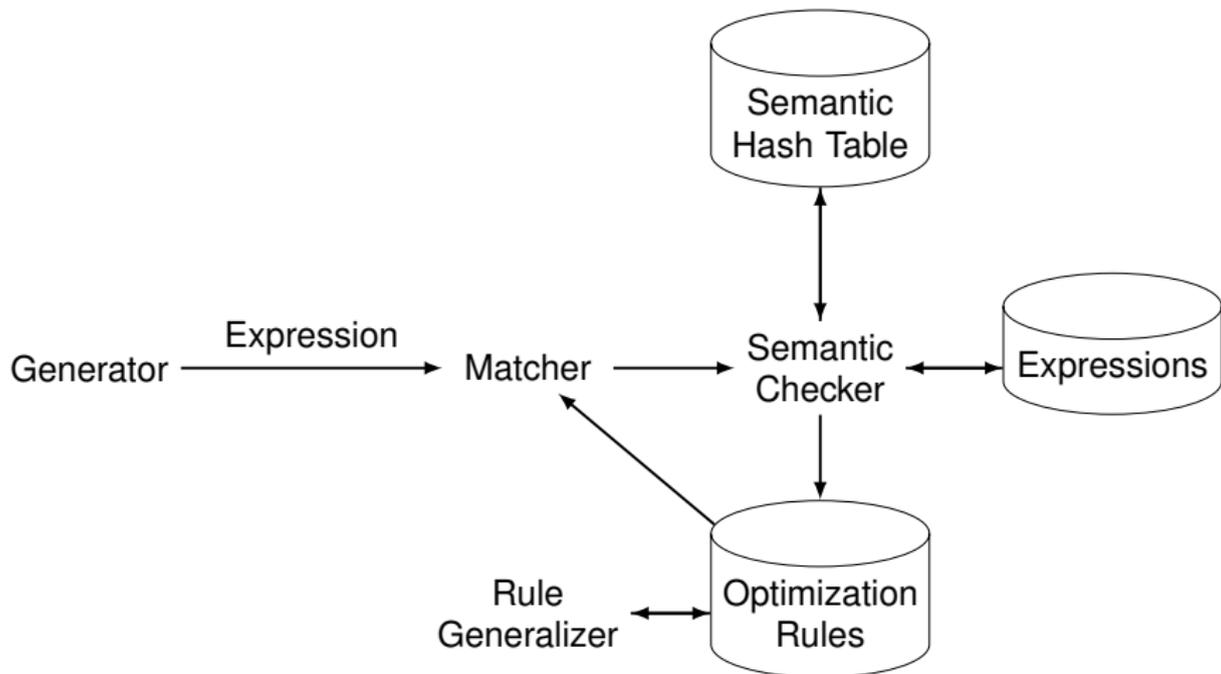
Example – Constant Folding Rules

Constant folding rules:

- $0 \ \& \ 0 \ \rightarrow \ 0$
 - $0 \ \& \ 1 \ \rightarrow \ 0$
 - $0 \ \& \ 2 \ \rightarrow \ 0$
 - ...
 - $255 \ \& \ 255 \ \rightarrow \ 255$
- } 2^{16} rules

Expected rule:

- $c0 \ \& \ c1 \ \rightarrow \ \text{eval}(c0 \ \& \ c1)$

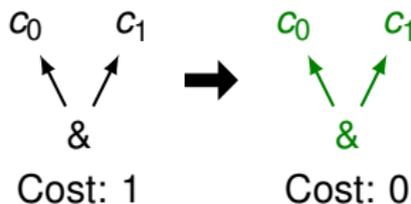


Example – Generalize Rules

Generalize constant folding rules:

1. Introduce *symbolic constants*

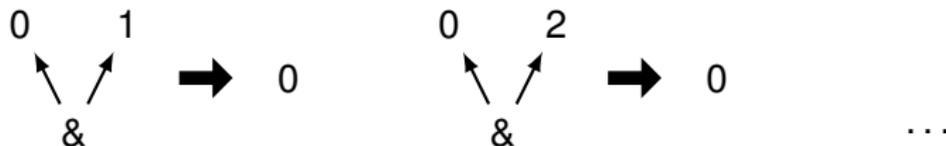
- Like variables
- Allow constant folding



Example – Generalize Rules

Generalize constant folding rules:

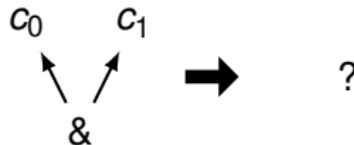
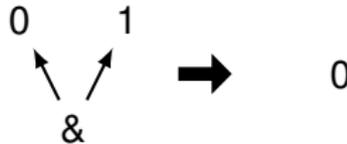
2. Collect syntactically equivalent rules



Example – Generalize Rules

Generalize constant folding rules:

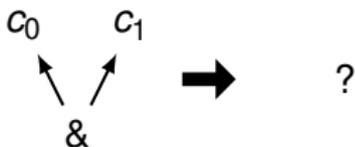
3. Replace constants of LHS with symbolic constants



Example – Generalize Rules

Generalize constant folding rules:

4. Iterate through generated expressions to find appropriate RHS



Example – Conditional Rules

Symbolic rules not sufficient:

- $(x \mid 2) \& 1 \rightarrow x \& 1$
- $(x \mid 1) \& 2 \rightarrow x \& 2$
- $(x \mid 1) \& 3 \not\rightarrow x \& 3$

Example – Conditional Rules

Symbolic rules not sufficient:

- $(x \mid 2) \& 1 \rightarrow x \& 1$
- $(x \mid 1) \& 2 \rightarrow x \& 2$
- $(x \mid 1) \& 3 \not\rightarrow x \& 3$

Solution:

- Conditional rule: $c0 \& c1 == 0 \Rightarrow (x \mid c0) \& c1 \rightarrow x \& c1$
- Iterate through generated expressions to find appropriate condition
 - Condition: $c0 \& c1 == 0$

Example – Result

OPTGEN finds 42 optimizations:

- 19 rules with symbolic constants
 - 8 rules with condition
 - 11 rules without condition
- 12 rules with non-symbolic constants
- 11 rules without constants

Example – Result

OPTGEN finds 42 optimizations:

- 19 rules with symbolic constants
 - 8 rules with condition
 - 11 rules without condition
- 12 rules with non-symbolic constants
- 11 rules without constants

Question

What happens if we use a bit width of 32 bit?

Example – Result

OPTGEN finds 42 optimizations:

- 19 rules with symbolic constants
 - 8 rules with condition
 - 11 rules without condition
- 12 rules with non-symbolic constants
- 11 rules without constants

Question

What happens if we use a bit width of 32 bit?



Extension to 32 Bit: Correctness

Basic idea:

- Generate rules for 8 bit
- Extend rules from 8 bit to 32 bit
- Verify extended rules for 32 bit

Extension of bit width:

- Rules without non-symbolic constants
 - Independent of bit width
 - $x \& x \rightarrow x$
- Rules with non-symbolic constants
 - Try to prepend or append 0/1 bits
 - $x \& 0xFF \rightarrow x$
 - $x \& 0xFF\ 000000 \rightarrow x$
 - $x \& 0xFF\ FFFFFFFF \rightarrow x$
 - $x \& 0x000000\ FF \rightarrow x$
 - $x \& 0xFFFFFFFF\ FF \rightarrow x$
 - Works fine in practice

Extension to 32 Bit: Completeness

Basic idea:

- Increase bit width until the number of rules stabilizes

Bit width	Number of rules
1	24
2	38
3	42
4	42
...	...
32	42

Drawback:

- Does not work for all operations

Full run:

- Operations: Constants, Minus, Not, Add, And, Or, Sub, Xor
- Cost limit: 2
- Generation: 8 bit
- Verification: 32 bit
- 6 h 7 min 0 s
- 1 046 568 kB

Testsuite:

- LLVM: 23 missing optimizations
- GCC: 27 missing optimizations
- ICC: 62 missing optimizations

	Optimization	Compiler		
		LLVM	GCC	ICC
2.	$-(x \& 0x80000000) \rightarrow x \& 0x80000000$	×	✓	×
6.	$(x \mid 0x80000000) + 0x80000000 \rightarrow x \& 0x7FFFFFFF$	✓	×	×
11.	$x \& (x + 0x80000000) \rightarrow x \& 0x7FFFFFFF$	✓	×	×
14.	$-x \& 1 \rightarrow x \& 1$	×	✓	×
17.	$x \mid (x + 0x80000000) \rightarrow x \mid 0x80000000$	✓	×	×
20.	$x \mid (x \oplus y) \rightarrow x \mid y$	✓	×	×
* 21.	$((c0 \mid -c0) \& \sim c1) == 0 \Rightarrow (x + c0) \mid c1 \rightarrow x \mid c1$	✓	×	✓
25.	$0 - (x \& 0x80000000) \rightarrow x \& 0x80000000$	×	✓	×
30.	$x \oplus (x + 0x80000000) \rightarrow 0x80000000$	✓	×	×
35.	$(0x7FFFFFFF - x) \oplus 0x80000000 \rightarrow \sim x$	×	✓	×
36.	$(0x80000000 - x) \oplus 0x80000000 \rightarrow -x$	×	✓	×
43.	$\sim(x + c) \rightarrow \sim c - x$	✓	×	×
54.	$\sim(c - x) \rightarrow x + \sim c$	✓	×	×
60.	$(c0 \& \sim c1) == 0 \Rightarrow (x \oplus c0) \mid c1 \rightarrow x \mid c1$	✓	×	×
Missing optimizations		5	9	13 (+ 32)

Unsupported Optimizations

	Optimization	Compiler		
		LLVM	GCC	ICC
5.	$x + (x \& 0x80000000) \rightarrow x \& 0x7FFFFFFF$	×	×	×
13.	$x \& (0x7FFFFFFF - x) \rightarrow x \& 0x80000000$	×	×	×
* 16.	$is_power_of_2(c1) \&\& c0 \& (2 * c1 - 1) == c1 - 1$ $\Rightarrow (c0 - x) \& c1 \rightarrow x \& c1$	×	×	×
19.	$x (0x7FFFFFFF - x) \rightarrow x 0x7FFFFFFF$	×	×	×
* 22.	$is_power_of_2(\sim c1) \&\& c0 \& (2 * \sim c1 - 1) == \sim c1 - 1$ $\Rightarrow (c0 - x) c1 \rightarrow x c1$	×	×	×
23.	$-x 0xFFFFFFFFE \rightarrow x 0xFFFFFFFFE$	×	×	×
26.	$0x7FFFFFFF - (x \& 0x80000000) \rightarrow x 0x7FFFFFFF$	×	×	×
27.	$0x7FFFFFFF - (x 0x7FFFFFFF) \rightarrow x \& 0x80000000$	×	×	×
28.	$0xFFFFFFFFE - (x 0x7FFFFFFF) \rightarrow x 0x7FFFFFFF$	×	×	×
29.	$(x \& 0x7FFFFFFF) - x \rightarrow x \& 0x80000000$	×	×	×
31.	$x \oplus (0x7FFFFFFF - x) \rightarrow 0x7FFFFFFF$	×	×	×
32.	$(x + 0x7FFFFFFF) \oplus 0x7FFFFFFF \rightarrow -x$	×	×	×
34.	$-x \oplus 0x80000000 \rightarrow 0x80000000 - x$	×	×	×
39.	$(0x7FFFFFFF - x) \oplus 0x7FFFFFFF \rightarrow x$	×	×	×
48.	$-x \oplus 0x7FFFFFFF \rightarrow x + 0x7FFFFFFF$	×	×	×
52.	$(x c) - c \rightarrow x \& \sim c$	×	×	×
57.	$-c0 == c1 \Rightarrow (x c0) + c1 \rightarrow x \& \sim c1$	×	×	×
62.	$0x7FFFFFFF - (x \oplus c) \rightarrow x \oplus (0x7FFFFFFF - c)$	×	×	×

OPTGEN

- is the first generator that supports arbitrary constants
- guarantees correctness and completeness of generated optimizations
- has revealed missing optimizations in all state-of-the-art compilers

There is more wisdom in the paper.

No

Optimization	Compiler		
	LLVM	GCC	ICC
1. $\sim\sim x \rightarrow x + 1$	✓	✓	×
2. $-(x \& 0x80000000) \rightarrow x \& 0x80000000$	×	✓	×
3. $\sim -x \rightarrow x - 1$	✓	✓	×
4. $x + \sim x \rightarrow 0xFFFFFFFF$	✓	✓	×
5. $x + (x \& 0x80000000) \rightarrow x \& 0x7FFFFFFF$	×	×	×
6. $(x 0x80000000) + 0x80000000 \rightarrow x \& 0x7FFFFFFF$	✓	×	×
7. $(x \& 0x7FFFFFFF) + (x \& 0x7FFFFFFF) \rightarrow x + x$	✓	✓	×
8. $(x \& 0x80000000) + (x \& 0x80000000) \rightarrow 0$	✓	✓	×
9. $(x 0x7FFFFFFF) + (x 0x7FFFFFFF) \rightarrow 0xFFFFFFFFE$	✓	✓	×
10. $(x 0x80000000) + (x 0x80000000) \rightarrow x + x$	✓	✓	×
11. $x \& (x + 0x80000000) \rightarrow x \& 0x7FFFFFFF$	✓	×	×
12. $x \& (x y) \rightarrow x$	✓	✓	×
13. $x \& (0x7FFFFFFF - x) \rightarrow x \& 0x80000000$	×	×	×
14. $-x \& 1 \rightarrow x \& 1$	×	✓	×
15. $(x + x) \& 1 \rightarrow 0$	✓	✓	×
16. $is_power_of_2(c1) \&\& c0 \& (2 * c1 - 1) == c1 - 1$ $\Rightarrow (c0 - x) \& c1 \rightarrow x \& c1$	×	×	×
Sum	23	27	62

Optimization	Compiler		
	LLVM	GCC	ICC
17. $x \mid (x + 0x80000000) \rightarrow x \mid 0x80000000$	✓	×	×
18. $x \mid (x \& y) \rightarrow x$	✓	✓	×
19. $x \mid (0x7FFFFFFF - x) \rightarrow x \mid 0x7FFFFFFF$	×	×	×
20. $x \mid (x \oplus y) \rightarrow x \mid y$	✓	×	×
21. $((c0 \mid -c0) \& \sim c1) == 0 \Rightarrow (x + c0) \mid c1 \rightarrow x \mid c1$	✓	×	✓
22. $is_power_of_2(\sim c1) \&\& c0 \& (2 * \sim c1 - 1) == \sim c1 - 1$ $\Rightarrow (c0 - x) \mid c1 \rightarrow x \mid c1$	×	×	×
23. $-x \mid 0xFFFFFFFFE \rightarrow x \mid 0xFFFFFFFFE$	×	×	×
24. $(x + x) \mid 0xFFFFFFFFE \rightarrow 0xFFFFFFFFE$	✓	✓	×
25. $0 - (x \& 0x80000000) \rightarrow x \& 0x80000000$	×	✓	×
26. $0x7FFFFFFF - (x \& 0x80000000) \rightarrow x \mid 0x7FFFFFFF$	×	×	×
27. $0x7FFFFFFF - (x \mid 0x7FFFFFFF) \rightarrow x \& 0x80000000$	×	×	×
28. $0xFFFFFFFFE - (x \mid 0x7FFFFFFF) \rightarrow x \mid 0x7FFFFFFF$	×	×	×
29. $(x \& 0x7FFFFFFF) - x \rightarrow x \& 0x80000000$	×	×	×
30. $x \oplus (x + 0x80000000) \rightarrow 0x80000000$	✓	×	×
31. $x \oplus (0x7FFFFFFF - x) \rightarrow 0x7FFFFFFF$	×	×	×
32. $(x + 0x7FFFFFFF) \oplus 0x7FFFFFFF \rightarrow -x$	×	×	×
Sum	23	27	62

Optimization	Compiler		
	LLVM	GCC	ICC
33. $(x + 0x80000000) \oplus 0x7FFFFFFF \rightarrow \sim x$	✓	✓	×
34. $-x \oplus 0x80000000 \rightarrow 0x80000000 - x$	×	×	×
35. $(0x7FFFFFFF - x) \oplus 0x80000000 \rightarrow \sim x$	×	✓	×
36. $(0x80000000 - x) \oplus 0x80000000 \rightarrow -x$	×	✓	×
37. $(x + 0xFFFFFFFF) \oplus 0xFFFFFFFF \rightarrow -x$	✓	✓	×
38. $(x + 0x80000000) \oplus 0x80000000 \rightarrow x$	✓	✓	×
39. $(0x7FFFFFFF - x) \oplus 0x7FFFFFFF \rightarrow x$	×	×	×
40. $x - (x \& c) \rightarrow x \& \sim c$	✓	✓	×
41. $x \oplus (x \& c) \rightarrow x \& \sim c$	✓	✓	×
42. $\sim x + c \rightarrow (c - 1) - x$	✓	✓	×
43. $\sim(x + c) \rightarrow \sim c - x$	✓	×	×
44. $-(x + c) \rightarrow -c - x$	✓	✓	×
45. $c - \sim x \rightarrow x + (c + 1)$	✓	✓	×
46. $\sim x \oplus c \rightarrow x \oplus \sim c$	✓	✓	×
47. $\sim x - c \rightarrow \sim c - x$	✓	✓	×
48. $-x \oplus 0x7FFFFFFF \rightarrow x + 0x7FFFFFFF$	×	×	×
Sum	23	27	62

	Optimization	Compiler		
		LLVM	GCC	ICC
49.	$-x \oplus 0xFFFFFFFF \rightarrow x - 1$	✓	✓	×
50.	$x \& (x \oplus c) \rightarrow x \& \sim c$	✓	✓	×
51.	$-x - c \rightarrow -c - x$	✓	✓	×
52.	$(x c) - c \rightarrow x \& \sim c$	×	×	×
53.	$(x c) \oplus c \rightarrow x \& \sim c$	✓	✓	×
54.	$\sim(c - x) \rightarrow x + \sim c$	✓	×	×
55.	$\sim(x \oplus c) \rightarrow x \oplus \sim c$	✓	✓	×
56.	$\sim c0 == c1 \Rightarrow (x \& c0) \oplus c1 \rightarrow x c1$	✓	✓	×
57.	$-c0 == c1 \Rightarrow (x c0) + c1 \rightarrow x \& \sim c1$	×	×	×
58.	$(x \oplus c) + 0x80000000 \rightarrow x \oplus (c + 0x80000000)$	✓	✓	×
59.	$((c0 -c0) \& c1) == 0 \Rightarrow (x \oplus c0) \& c1 \rightarrow x \& c1$	✓	✓	×
60.	$(c0 \& \sim c1) == 0 \Rightarrow (x \oplus c0) c1 \rightarrow x c1$	✓	×	×
61.	$(x \oplus c) - 0x80000000 \rightarrow x \oplus (c + 0x80000000)$	✓	✓	×
62.	$0x7FFFFFFF - (x \oplus c) \rightarrow x \oplus (0x7FFFFFFF - c)$	×	×	×
63.	$0xFFFFFFFF - (x \oplus c) \rightarrow x \oplus (0xFFFFFFFF - c)$	✓	✓	×
Sum		23	27	62

Optimization	Compiler		
	LLVM	GCC	ICC
1. $\sim(x \mid \sim y) \rightarrow \sim x \& y$	×	✓	
2. $\sim(x \& \sim y) \rightarrow \sim x \mid y$	×	✓	
3. $(x+x) \& (y+y) \rightarrow (x \& y) + (x \& y)$	×		
4. $(x+x) \mid (y+y) \rightarrow (x \mid y) + (x \mid y)$	×		
5. $(x \& y) \mid (z \& y) \rightarrow y \& (x \mid z)$	✓	×	✓
6. $x - ((x - y) + (x - y)) \rightarrow y + (y - x)$		✓	×
7. $(x - y) - (x + z) \rightarrow -(y + z)$	✓	✓	×
8. $((x - y) + (x - y)) - x \rightarrow x - (y + y)$	✓	✓	×
9. $(x+x) \oplus (y+y) \rightarrow (x \oplus y) + (x \oplus y)$	×		
10. $(x \& y) \oplus (z \& y) \rightarrow y \& (x \oplus z)$	✓	×	✓

State-of-the-art compilers apply optimization rules even if the operands are shared. If the compiler supports the optimization ✓/× indicates whether the compiler prevents the optimization in case of shared operands. If the compiler does not support the optimization the item is left blank.